

HEURISTIC TWO-LEVEL LOGIC OPTIMIZATION

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Outline

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- Heuristic logic minimization.
- Principles.
- Operators on logic covers.
- Espresso.

Heuristic minimization

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- Provide irredundant covers with 'reasonably small' cardinality.
- Fast and applicable to many functions.
- Avoid bottlenecks of exact minimization:
 - Prime generation and storage.
 - Covering.

Heuristic minimization Principles

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- Local minimum cover:
 - Given initial cover.
 - Make it prime.
 - Make it irredundant.
- Iterative improvement:
 - Improve on cardinality by 'modifying' the implicants.

Heuristic minimization Operators

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- **Expand:**
 - Make implicants prime.
 - Remove covered implicants.
- **Reduce:**
 - Reduce size of each implicant while preserving cover.
- **Reshape:**
 - Modify implicant pairs: enlarge one and reduce the other.
- **Irredundant:**
 - Make cover irredundant.

Example

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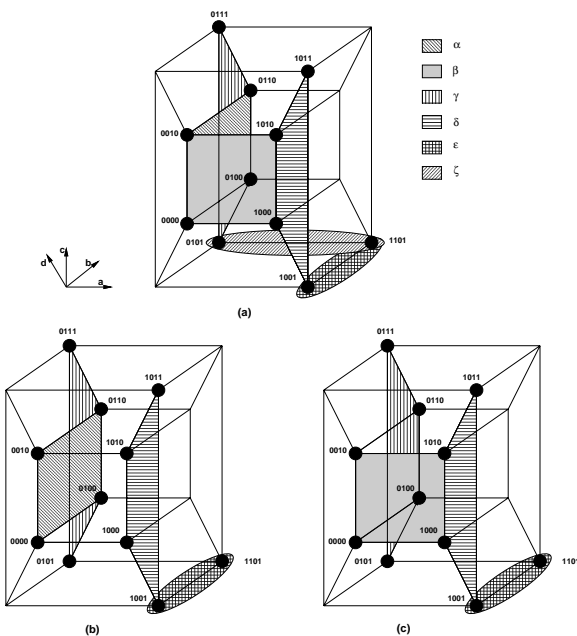
```

0000  1
0010  1
0100  1
0110  1
1000  1
1010  1
0101  1
0111  1
1001  1
1011  1
1101  1
    
```

α	0**0	1
β	*0*0	1
γ	01**	1
δ	10**	1
ϵ	1*01	1
ζ	*101	1

Example

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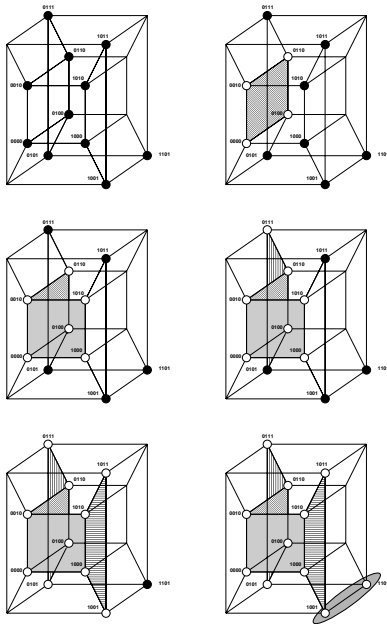
Example Expansion

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- Expand 0000 to $\alpha = 0**0$.
 - Drop 0100, 0010, 0110 from the cover.
- Expand 1000 to $\beta = *0*0$.
 - Drop 1010 from the cover.
- Expand 0101 to $\gamma = 01**$.
 - Drop 0111 from the cover.
- Expand 1001 to $\delta = 10**$.
 - Drop 1011 from the cover.
- Expand 1101 to $\epsilon = 1*01$.
- Cover is: $\{\alpha, \beta, \gamma, \delta, \epsilon\}$.

Example

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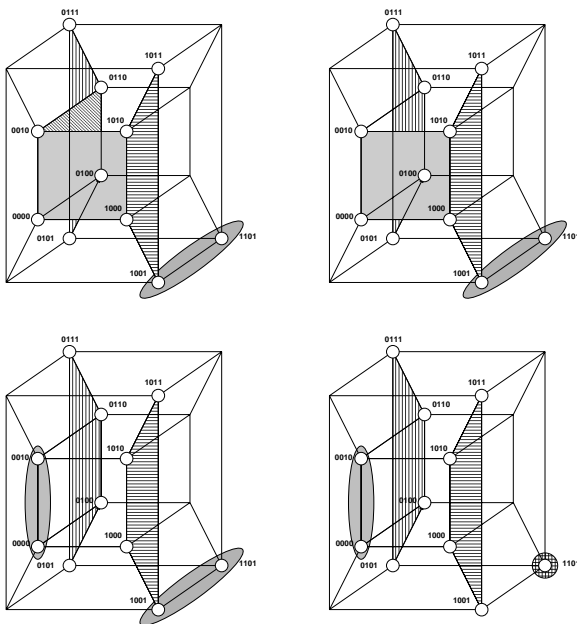
Example Reduction

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- Reduce $0**0$ to nothing.
- Reduce $\beta = *0*0$ to $\tilde{\beta} = 00*0$
- Reduce $\epsilon = 1*01$ to $\tilde{\epsilon} = 1101$
- Cover is: $\{\tilde{\beta}, \gamma, \delta, \tilde{\epsilon}\}$.

Example

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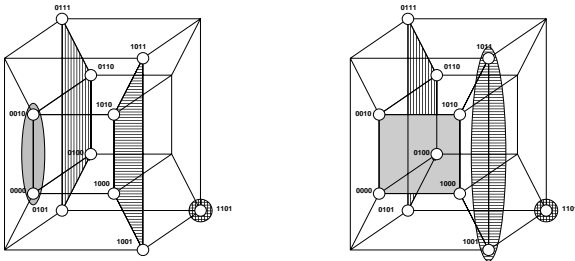
Example Reshape

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- Reshape $\{\tilde{\beta}, \delta\}$ to: $\{\beta, \tilde{\delta}\}$
 – where $\tilde{\delta} = 10*1$.
- Cover is: $\{\beta, \gamma, \tilde{\delta}, \tilde{\epsilon}\}$.

Example

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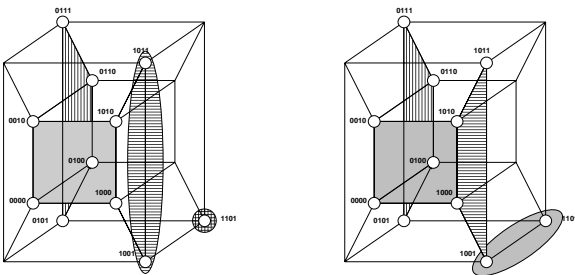
Example Second expansion

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- Expand $\tilde{\delta} = 10 * 1$ to $\delta = 10 * *$.
- Expand $\tilde{\epsilon} = 1101$ to $\epsilon = 1 * 01$.

Example

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Example (MINI summary)

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- Expansion:
 - Cover: $\{\alpha, \beta, \gamma, \delta, \epsilon\}$.
 - Prime, redundant, minimal w.r. to scc.
- Reduction:
 - α eliminated.
 - $\beta = *0*0$ reduced to $\tilde{\beta} = 00*0$.
 - $\epsilon = 1*01$ reduced to: $\tilde{\epsilon} = 1101$.
 - Cover: $\{\tilde{\beta}, \gamma, \delta, \tilde{\epsilon}\}$.
- Reshape:
 - $\{\tilde{\beta}, \delta\}$ reshaped to: $\{\beta, \tilde{\delta}\}$ where $\tilde{\delta} = 10*1$.
- Second expansion:
 - Cover: $\{\beta, \gamma, \delta, \epsilon\}$.
 - Prime, irredundant.

Alternative example (ESPRESSO)

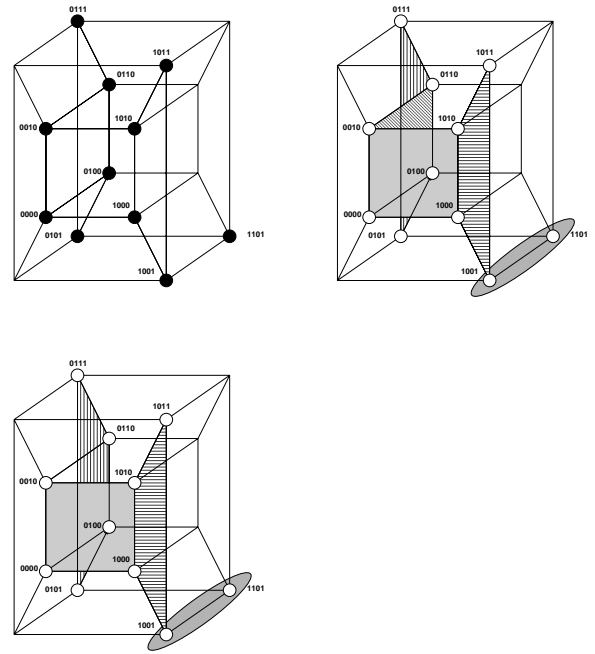
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- Expansion:
 - Cover: $\{\alpha, \beta, \gamma, \delta, \epsilon\}$.
 - Prime, redundant, minimal w.r. to scc.

- Irredundant:
 - Cover: $\{\beta, \gamma, \delta, \epsilon\}$.
 - Prime, irredundant.

Example

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Expand naive implementation

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- For each implicant
 - For each *care* literal
 - * Raise it to *don't care* if possible.
 - Remove all covered implicants.

- Problems:
 - Validity check.
 - Order of expansions.

Validity check

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- Espresso, MINI:
 - Check *intersection* of expanded implicant with OFF-set.
 - Requires complementation.

- Presto:
 - Check *inclusion* of expanded implicant in the union of the ON-set and DC-set.
 - Can be reduced to recursive tautology check.

Heuristics

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- Expand first cubes that are unlikely to be covered by other cubes.
- Selection:
 - Compute vector of column sums.
 - *Weight*: inner product of cube and vector.
 - Sort implicants in ascending order of weight.
- Rationale:
 - Low weight correlates to having few 1s in densely populated columns.

Example

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- $f = a'b'c' + ab'c' + a'bc' + a'b'c$
DC-set = abc'

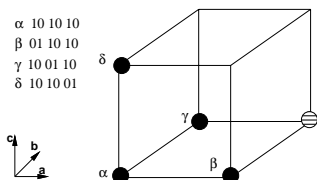
10	10	10
01	10	10
10	01	10
10	10	01

- Ordering:
 - Vector: $[313131]^T$
 - Weights: (9, 7, 7, 7).
- Select second implicant.

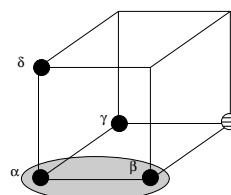
Example (2)

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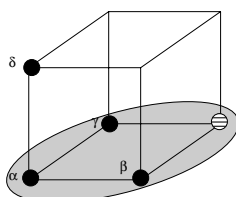
α 10 10 10
 β 01 10 10
 γ 10 01 10
 δ 10 10 01



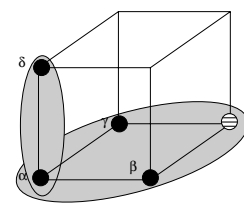
(a)



(b)



(c)



(d)

Example (3)

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- OFF-set:

01	11	01
11	01	01

- Expand 01 10 10:
 - 11 10 10 valid.
 - 11 11 10 valid.
 - 11 11 11 invalid.

- Update cover to:

11	11	10
10	10	01

Example (4)

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11 11 10
10 10 01

- Expand 10 10 01:
 - 11 10 01 invalid.
 - 10 11 01 invalid.
 - 10 10 11 valid.

- Expanded cover:

11 11 10
10 10 11

Definitions

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- *free*:
 - Set of entries that can be raised to 1.
- *Overexpanded cube*
 - Cube whose entries in *free* are raised.
- *Feasible cover*
 - Expand a cube to cover another one while keeping it as an implicant of the function.

Expand

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- Smarter heuristics for choosing literals to be expanded.
- Four step procedure in Espresso.
- Rationale:
 - Raise literals so that expanded implicant:
 - * Covers a maximal set of cubes.
 - * Making it as large as possible.

Expand in ESPRESSO

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- *Determine the essential parts.*
 - Determine which entries can never be raised, and remove them from *free*.
 - Determine which parts can always be raised, raise them, and remove them from *free*.
- *Detection of feasibly covered cubes.*
 - If there is an implicant β whose supercube with α is feasible, repeat the following steps.
 - * Raise the appropriate entry of α and remove it from *free*.
 - * Remove from *free* entries that can never be raised or that can always be raised and update α .
- *Expansion guided by the overexpanded cube.*
 - While the overexpanded cube of α covers some other cubes of F , repeat the following steps.
 - * Raise a single entry of α as to overlap a maximum number of those cubes.
 - * Remove from *free* entries that can never be raised or that can always be raised and update α .
- *Find the largest prime implicant.*
 - Formulate a covering problem and solve it by a heuristic method.

Reduce

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- Sort implicants:
 - Heuristic: sort by descending weight.
- For each implicant:
- Lower as many * as possible to 1 or 0.
- *Theorem*:
 - Let $\alpha \in F$ and $Q = F \cup D - \{\alpha\}$.
Then, the maximally reduced cube is:
 $\tilde{\alpha} = \alpha \cap \text{supercube}(Q'_\alpha)$.

Example

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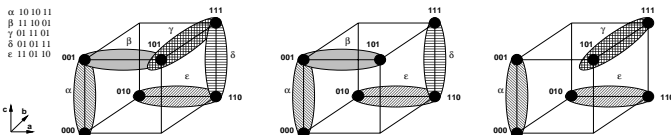
- Expanded cover:

$$\begin{array}{ccc} 11 & 11 & 10 \\ 10 & 10 & 11 \end{array}$$
- Select first implicant:
 - cannot be reduced.
- Select second implicant:
 - Reduced to 10 10 01
- Reduced cover:

$$\begin{array}{ccc} 11 & 11 & 10 \\ 10 & 10 & 01 \end{array}$$

Irredundant cover

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Irredundant cover

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- *Relatively essential set* E^r
 - Implicants covering some minterms of the function not covered by other implicants.
- *Totally redundant set* R^t
 - Implicants covered by the relatively essentials.
- *Partially redundant set* R^p
 - Remaining implicants.

Irredundant cover

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- Find a subset of R^p that, together with E^r , covers the function.
- Modification of the tautology algorithm:
 - Each cube in R^p is covered by other cubes.
 - Find mutual covering relations.
- Reduces to a covering problem:
 - Heuristic algorithm.

Example

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α	10	10	11
β	11	10	01
γ	01	11	01
δ	01	01	11
ϵ	11	01	10

- $E^r = \{\alpha, \epsilon\}$
- $R^t = \emptyset$
- $R^p = \{\beta, \gamma, \delta\}$.

Example (2)

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- Covering relations:
 - β is covered by $\{\alpha, \gamma\}$.
 - γ is covered by $\{\beta, \delta\}$.
 - δ is covered by $\{\gamma, \epsilon\}$.
- Minimum cover: $\gamma \cup E^r$

Espresso algorithm

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- Compute the complement.
- Extract essentials.
- Iterate:
 - *Expand, irredundant, reduce.*
- Cost functions:
 - Cover cardinality ϕ_1 .
 - Weighed sum of cube and literal count ϕ_2 .

Espresso algorithm

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```
espresso(F, D){
  R = complement(F ∪ D);
  F = expand(F, R);
  F = irredundant(F, D);
  E = essentials(F, D);
  F = F - E;
  D = D ∪ E;
  repeat {
    φ2 = cost(F);
    repeat {
      φ1 = |F|;
      F = reduce(F, D);
      F = expand(F, R);
      F = irredundant(F, D);
    } until ( |F| ≥ φ1);
    F = last_gasp(F, D, R);
  } until ( cost(F) ≥ φ2);
  F = F ∪ E;
  D = D - E;
  F = make_sparse(F, D, R);
}
```

Summary heuristic minimization

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- Heuristic minimization is iterative.
- Few operators applied to covers.
- Underlying mechanism:
 - Cube operation.
 - Unate recursive paradigm.
- Efficient algorithms.