LIBRARY BINDING

©*Giovanni De Micheli*

Stanford University

Outline

- Modeling and problem analysis.
- Rule-based systems for library binding.
- Algorithms for library binding:
  - *Structural covering/matching.*
  - *Boolean covering/matching.*
- Concurrent optimization and binding.

Library binding

- Given an unbound logic network and a set of library cells:
  - Transform into an interconnection of instances of library cells.
  - Optimize *area*, (under *delay* constraints.)
  - Optimize *delay*, (under *area* constraints.)
  - Optimize *power*, (under *delay* constraints.)
- Called also *technology mapping*:
  - Method used for re-designing circuits in different technologies.

Library models

- Combinational elements:
  - Single-output functions:
    - e.g. AND, OR, AOI.
  - Compound cells: e.g. adders, encoders.
- Sequential elements:
  - Registers, counters.
- Miscellaneous:
  - Schmitt triggers.

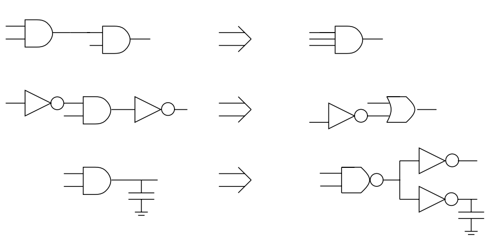
Major approaches

- Rule-based systems:
  - Mimic designer activity.
  - Handle all types of cells.
- Heuristic algorithms:
  - Restricted to single-output combinational cells.
- Most tools use a combination of both.

Rule-based library binding

- Binding by stepwise transformations.
- Data-base:
  - Set of patterns associated with best implementation.
- Rules:
  - Select subnetwork to be mapped.
  - Handle high-fanout problems, buffering, etc.

Example



Strategies

- Search for a sequence of transformations.
- Search space:
  - *Breadth* (options at each step).
  - *Depth* (look-ahead).
- *Meta-rules* determine dynamically breadth and depth.

Rule-based library binding

© GDM

- Advantages:
  - Applicable to all kinds of libraries.
- Disadvantages:
  - Large rule data-base:
    - Completeness issue.
    - Formal properties of bound network.
  - Data-base updates.

Algorithms for library binding

© GDM

- Mainly for single-output combinational cells.
- Fast and efficient:
  - Quality comparable to rule-based systems.
- Library description/update is simple:
  - Each cell modeled by its function or equivalent pattern.

Problem analysis

© GDM

- Matching:
  - A cell matches a sub-network if their terminal behavior is the same.
  - Input-variable *assignment* problem.
- Covering:
  - A cover of an unbound network is a partition into subnetworks which can be replaced by library cells.

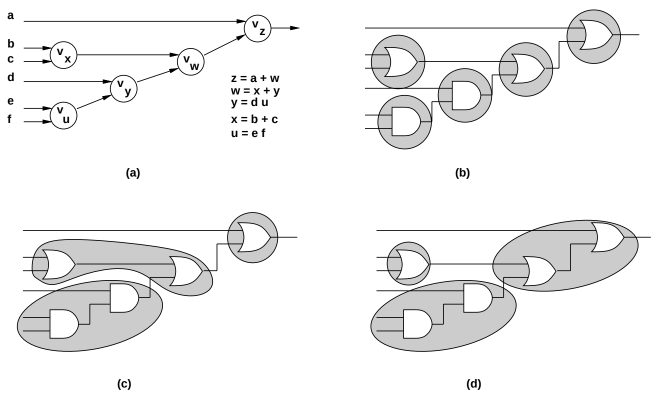
Assumptions

© GDM

- Network granularity is fine.
  - Decomposition into *base* functions.
    - 2-input $AND, OR, NAND, NOR$.
- Trivial binding:
  - Replacement of each vertex by base cell.

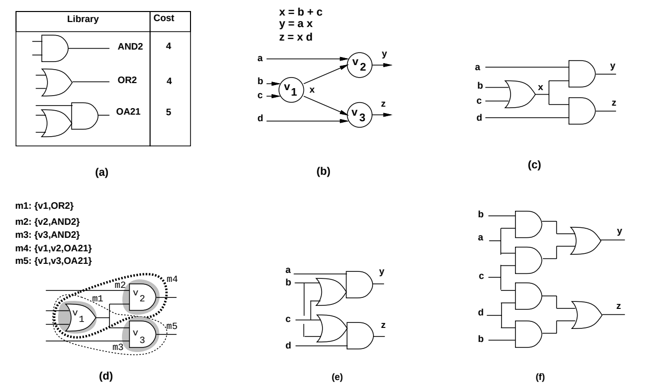
Example

© GDM



Example

© GDM



Example

© GDM

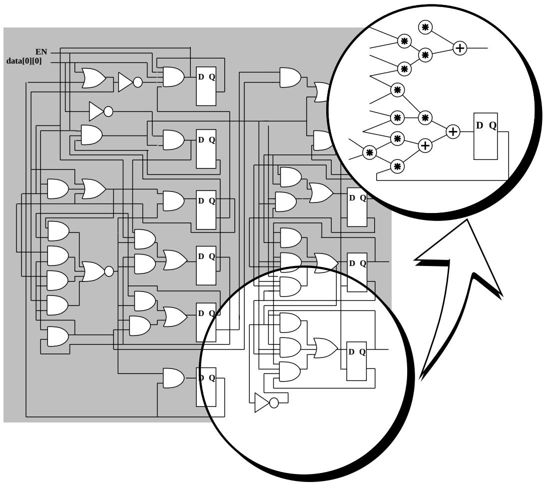
- Vertex covering:
  - Covering v_1: $(m_1 + m_4 + m_5)$.
  - Covering v_2: $(m_2 + m_4)$.
  - Covering v_3: $(m_3 + m_5)$.
- Input compatibility:
  - Match m_2 requires m_1:
    * $(m_2' + m_1)$.
  - Match m_3 requires m_1:
    * $(m_3' + m_1)$.
- Overall *binate* clause:
  - $(m_1 + m_4 + m_5)(m_2 + m_4)(m_3 + m_5)(m_2' + m_1)(m_3' + m_1) = 1$

Heuristic algorithms

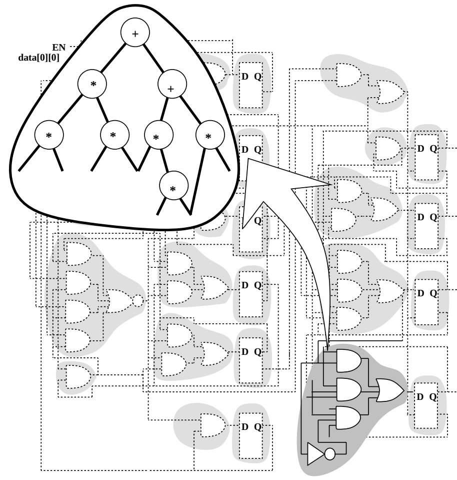
© GDM

- Decomposition:
  - Cast network and library in standard form.
  - Decompose into *base functions*.
  - Example: NAND2 and INV.
- Partitioning:
  - Break network into *cones*.
  - Reduce to many multi-input single-output subnetworks.
- Covering:
  - Cover each subnetwork by library cells.

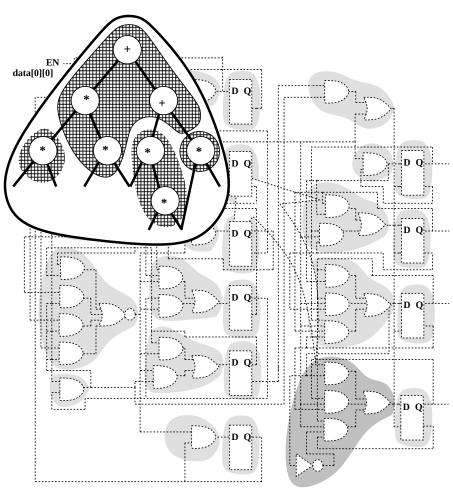
Decomposition



Partitioning



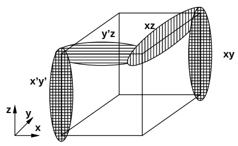
Covering



Heuristic algorithms

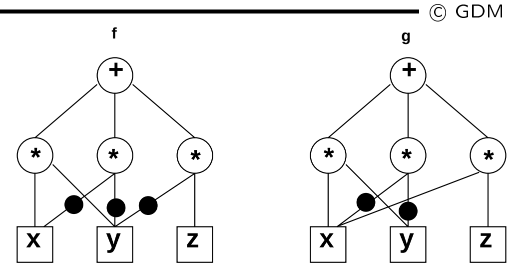
- Structural approach:
  - Model functions by *patterns*.
    - ∗ Example: trees, dags.
  - Rely on *pattern matching* techniques.
- Boolean approach:
  - Use Boolean models.
  - Solve *tautology* problem.
  - More powerful.

Example
Boolean versus structural matching



- $f = xy + x'y' + y'z$
- $g = xy + x'y' + xz$
- Function equality is a tautology:
  - Boolean match.
- Patterns may be different:
  - Structural match may not be found.

Example
Boolean versus structural matching

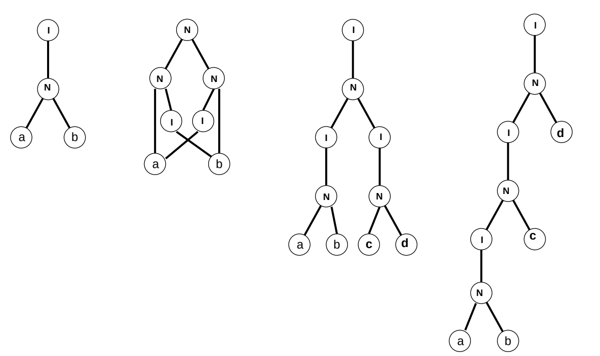


- $f = xy + x'y' + y'z$
- $g = xy + x'y' + xz$
- Patterns do not match.

Structural matching and covering

- Expression patterns:
  - Represented by dags.
- Identify pattern dags in network:
  - Sub-graph isomorphism.
- Simplification:
  - Use tree patterns.

Example

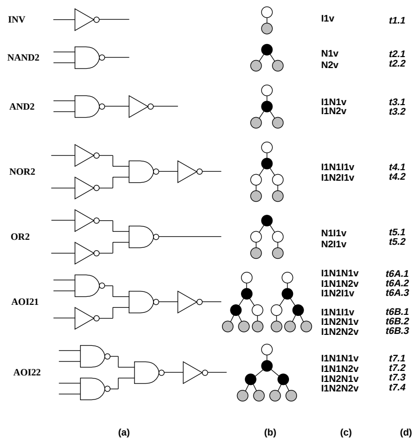


(a) (b) (c) (d)

Tree-based matching

- Network:
  - Partitioned and decomposed:
    - NOR2 (or NAND2) + INV.
    - Generic base functions.
  - *Subject tree.*
- Library:
  - Represented by trees.
  - Possibly more than one tree per cell.
- Pattern recognition:
  - Simple binary tree match.
  - Aho-Corasick automaton.

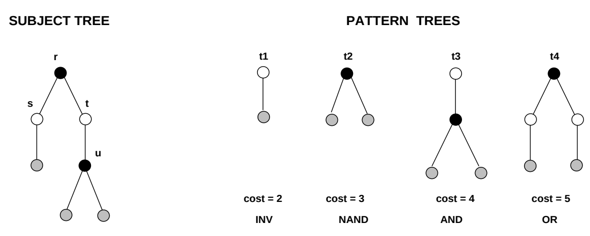
Simple library



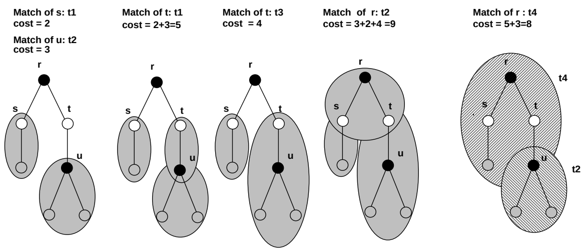
Tree covering

- Dynamic programming:
  - Visit subject tree bottom-up.
- At each vertex:
  - Attempt to match:
    - Locally rooted subtree.
    - All library cells.
- *Optimum* solution, for the subtree.

Example



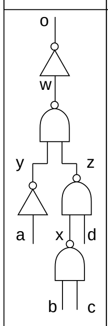
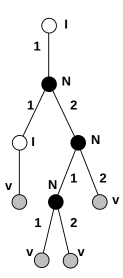
Example



Example

- Minimum-area cover.
- Area costs:
  - INV:2; NAND2:3; AND2:4; AOI21:6.
- Best choice:
  - AOI21 fed by a NAND2 gate.

Example

Network	Subject graph	Vertex	Match	Gate	Cost
		x	t2	NAND2(b,c)	3
		y	t1	INV(a)	2
		z	t2	NAND2(x,d)	2* 3 = 6
		w	t2	NAND2(y,z)	3 * 3 + 2 = 11
		o	t1	INV(w)	3 * 3 + 2 * 2 = 13
			t3	AND2(y,z)	2 * 3 + 4 + 2 = 12
			t6B	AOI21(x,d,a)	3 + 6 = 9

Minimum delay cover

- Dynamic programming approach.
- Cost related to gate delay.
- Delay modeling:
  - Constant gate delay.
    - ∗ Straightforward.
  - Load-dependent delay:
    - ∗ Load fanout unknown.
    - ∗ Binning techniques.

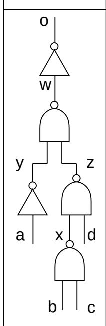
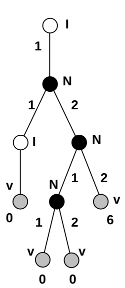
Minimum delay cover constant delays

- The cell pattern tree and the rooted subtree are isomorphic.
  - The vertex is labeled with the cell delay.
- The cell tree is isomorphic to a subtree with leaves L.
  - The vertex is labeled with the cell cost plus the *maximum* of the labels of L.

Example

- Inputs data-ready times are 0 except for $t_d = 6$.
- Constant delays:
  - INV:2; NAND2:4; AND2:5; AOI21:10.
- Compute *data-ready* times bottom-up:
  - $t_x = 4, t_y = 2; t_z = 10 t_w = 14$.
- Best choice:
  - AND2, two NAND2 and an INV gate.

Example

Network	Subject graph	Vertex	Match	Gate	Cost
		x	t2	NAND2(b,c)	4
		y	t1	INV(a)	2
		z	t2	NAND2(x,d)	6 + 4 = 10
		w	t2	NAND2(y,z)	10 + 4 = 14
		o	t1	INV(w)	14 + 2 = 16
			t3	AND2(y,z)	10 + 5 = 15
			t6B	AOI21(x,d,a)	10 + 6 = 16

Minimum delay cover load-dependent delays

- Model:
  - Assume a finite set of load values.
- Dynamic programming approach:
  - Compute an array of solutions for each possible load.
  - For each input to a matching cell the best match for any load is selected.
- *Optimum* solution, when all possible loads are considered.

Example

© GDM

- Inputs data-ready times are 0 except for $t_d = 6$.
- Load-dependent delays:
  - INV:1+l; NAND2:3+l; AND2:4+l; AOI21:9+l.
- Loads:
  - INV:1; NAND2:1; AND2:1; AOI21:1.
- Same solution as before.

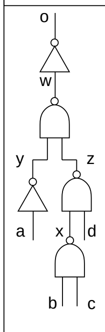
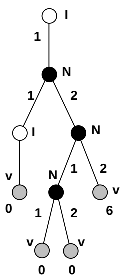
Example

© GDM

- Inputs data-ready times are 0 except for $t_d = 6$.
- Load-dependent delays:
  - INV:1+l; NAND2:3+l; AND2:4+l; AOI21:9+l; SINV:1+0.5l;.
- Loads:
  - INV:1; NAND2:1; AND2:1; AOI21:1; SINV:2.
- Assume output load is 1:
  - Same solution as before.
- Assume output load is 5:
  - Solution uses SINV cell.

Example

© GDM

Network	Subject graph	Vertex	Match	Gate	Cost Load=1	Cost Load=2	Cost Load=5
		x	*t2*	NAND2(b,c)	4	5	8
		y	*t1*	INV(a)	2	3	6
		z	*t2*	NAND2(x,d)	10	11	14
		w	*t2*	NAND2(y,z)	14	15	18
		o	*t1*	INV(w)			20
			t3	AND2(y,z)			19
			t6B	AOI21(x,d,a)			20
				SINV(w)			18.5

Library binding and polarity assignment

© GDM

- Search for lower cost solution by not constraining the signal polarities.
- Most circuit allow us to choose the input/output signal polarities.
- Approaches:
  - Structural covering.
  - Boolean covering.

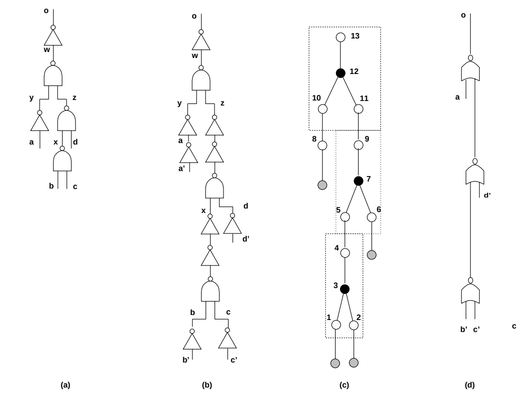
Structural covering and polarity assignment

© GDM

- Pre-process subject network:
  - Add inverter pairs between NANDs.
  - Provide signals with both polarity.
- Add inverter-pair cell to the library:
  - To eliminate unneeded pairs.
  - Cell corresponds to a connection with zero cost.

Example

© GDM



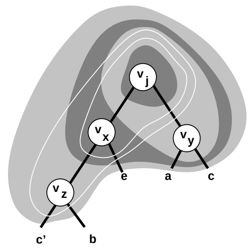
Boolean covering

© GDM

- Decompose network into base functions.
- When considering vertex v_i:
  - Construct *clusters* by local elimination.
  - Several functions associated with v_i.
- Limit size and depth of clusters.

Example

© GDM



$$
\begin{aligned}
f_{j,1} &= xy; \\
f_{j,2} &= x(a+c); \\
f_{j,3} &= (e+z)y; \\
f_{j,4} &= (e+z)(a+c); \\
f_{j,5} &= (e+c'+d)y; \\
f_{j,6} &= (e+c'+d)(a+c);
\end{aligned}
$$

Boolean matching $\mathcal{P}$-equivalence

- *Cluster function* $f(\mathbf{x})$: sub-network behavior.
- *Pattern function* $g(\mathbf{y})$: cell behavior.
- $\mathcal{P}$-equivalence:
  - Exists a permutation operator $\mathcal{P}$, such that $f(\mathbf{x}) = g(\mathcal{P}\ \mathbf{x})$ is a tautology?
- Approaches:
  - Tautology check over all input permutations.
  - Multi-rooted pattern ROBDD capturing all permutations.

Input/output polarity assignment

- Allow for reassignment of input/output polarity.
- $\mathcal{NPN}$classification of Boolean functions.
- $\mathcal{NPN}$-equivalence:
  - Exists a permutation matrix $\mathcal{P}$, and complementation operators $\mathcal{N}_i, \mathcal{N}_o$ such that $f(\mathbf{x}) = \mathcal{N}_o\ g(\mathcal{P}\ \mathcal{N}_i\ \mathbf{x})$ is a tautology?
- Variations:
  - $\mathcal{N}$-equivalence, $\mathcal{PN}$-equivalence

Boolean matching

- *Pin assignment* problem.
  - Map cluster variables $\mathbf{x}$ to pattern vars $\mathbf{y}$.
  - Characteristic equation: $\mathcal{A}(\mathbf{x}, \mathbf{y}) = 1$.
- Pattern function under variable assignment:
  - $g_{\mathcal{A}}(\mathbf{x}) = \mathcal{S}_{\mathbf{y}} \mathcal{A}(\mathbf{x}, \mathbf{y})\ g(\mathbf{y})$
- *Tautology problem.*
  - $f(\mathbf{x}) \overline{\oplus}\ g_{\mathcal{A}}(\mathbf{x})$
  - $\forall_{\mathbf{x}}(f(\mathbf{x}) \overline{\oplus}\ \mathcal{S}_{\mathbf{y}}\ (\mathcal{A}(\mathbf{x}, \mathbf{y})\ g(\mathbf{y})))$

Example

- Assign x_1 to y_2' and x_2 to y_1.
- Characteristic equation:
  - $A(x_1, x_2, y_1, y_2) = (x_1 \oplus y_2)(x_2 \overline{\oplus}\ y_1)$
- AND pattern function:
  - $g = y_1 y_2$
- Pattern function under assignment:
  - $\mathcal{S}_{y_1,y_2} \mathcal{A} g =$
    $= \mathcal{S}_{y_1,y_2} (x_1 \oplus y_2)(x_2 \overline{\oplus} y_1) y_1 y_2 = x_2 x_1'$

Signatures and filters

© GDM

- Capture some properties of Boolean functions.
- If signatures do not match, there is no match.
- Used as filters to reduce computation.
- Signatures:
  - Unateness.
  - Symmetries.
  - Co-factor sizes.
  - Spectra.

Filters based on unateness and symmetries

© GDM

- Any pin assignment must associate
  - unate (binate) variables in $f(\mathbf{x})$ with unate (binate) variables in $g(\mathbf{y})$.
- Variables or groups of variables
  - that are interchangeable in $f(\mathbf{x})$ must be interchangeable in $g(\mathbf{y})$.

Example

© GDM

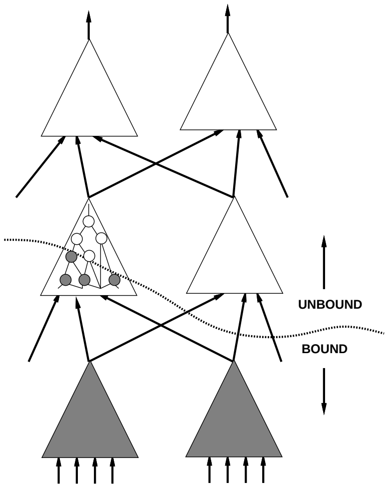
- Cluster function: $f = abc$.
  - Symmetries:$\{(a, b, c)\}$ – unate.
- Pattern functions:
  - $g_1 = a + b + c$
    - Symmetries:$\{(a, b, c)\}$ – unate.
  - $g_2 = ab + c$
    - Symmetries:$\{(a, b)(c)\}$ – unate.
  - $g_3 = abc' + a'b'c$
    - Symmetries:$\{(a, b, c)\}$ – binate.

Concurrent optimization and library binding

© GDM

- Motivation:
  - Logic simplification is usually done prior to binding.
  - Logic simplification/substitution can be combined with binding.
- Mechanism:
  - Binding induces some *don't care* conditions.
  - Exploit *don't cares* as degrees of freedom in matching.

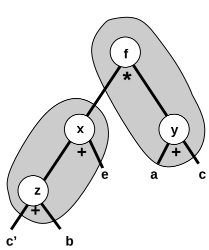
Example



Boolean matching with *don't care* conditions

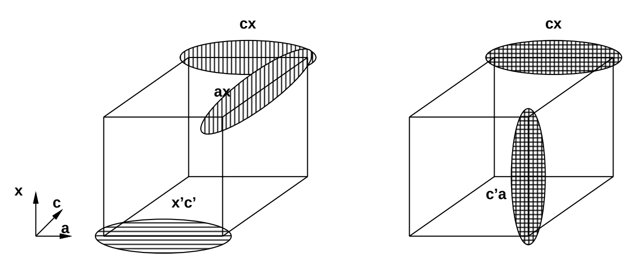
- Given $f(\mathbf{x}), f_{DC}(\mathbf{x})$ and $g(\mathbf{y})$:
  - g matches f if g is equivalent to $\tilde{f}$ where $f \cdot f'_{DC} \leq \tilde{f} \leq f + f_{DC}$
- Matching condition:
  - $\forall_{\mathbf{x}} (f_{DC}(\mathbf{x}) + f(\mathbf{x}) \overline{\oplus} \mathcal{S}_{\mathbf{y}} \, (\mathcal{A}(\mathbf{x}, \mathbf{y}) \, g(\mathbf{y})))$

Example

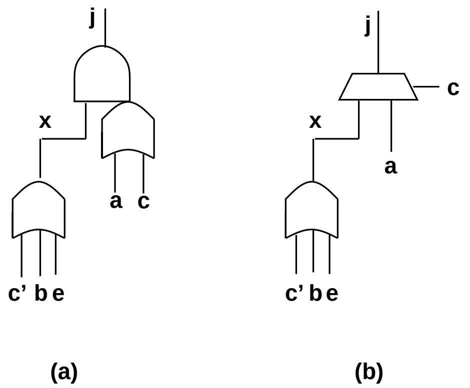


- Assume v_x is bound to $OR3(c', b, e)$.
- *Don't care* set includes $x \oplus (c' + b + e)$.
- Consider $f_j = x(a + c)$ with $CDC = x'c'$.
- No simplification. Mapping into AOI gate.
- Matching with DC. Mapping into MUX gate.

Example



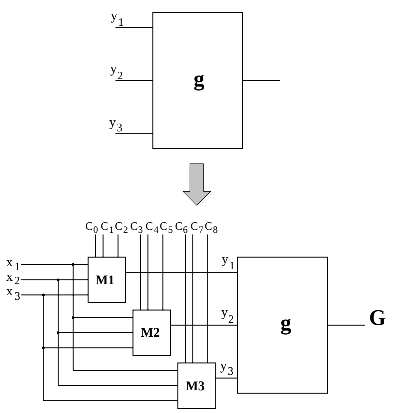
Example



Extended matching

- Augment pattern function with mux function.
  - Each cell input can be routed to any cluster input (or voltage rail).
  - Input polarity can be changed.
  - Cell and cluster may differ input size.
- Define composite function $G(\mathbf{x}, \mathbf{c})$:
  - Pin assignment is determining $\mathbf{c}$.
- Matching formula: $M(\mathbf{c}) = \forall_{\mathbf{x}} [G(\mathbf{x}, \mathbf{c}) \ \overline{\oplus} \ f(\mathbf{x})]$

Example

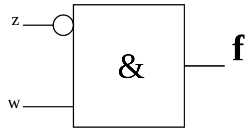


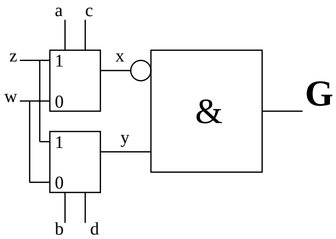
- $g = y_1 + y_2 \ y_3'$
- $y_1(\mathbf{c}, \mathbf{x}) = (c_0 c_1 x_1 + c_0 c_1' x_2 + c_0' c_1 x_3) \oplus c_2$
- $G = y_1(\mathbf{c}, \mathbf{x}) + y_2(\mathbf{c}, \mathbf{x}) \ y_3(\mathbf{c}, \mathbf{x})'$

Extended matching modeling

- Model composite functions by ROBDDs.
  - Assume: n-input cluster and m-input cell.
  - For each cell input:
    * $\lceil log_2 n \rceil$ variables for pin permutation.
    * One variable for input polarity.
  - Total size of $\mathbf{c}$: $m(\lceil log_2 n \rceil + 1)$.
- A match exists if there is at least one value of $\mathbf{c}$ satisfying $M(\mathbf{c}) = \forall_{\mathbf{x}} [G(\mathbf{x}, \mathbf{c}) \ \overline{\oplus} \ f(\mathbf{x})]$.

Example





- $g = x'y, \; f = wz'$
- $G(a,b,c,d,w,z) = (c \oplus (za + wa'))'(d \oplus (zb + wb'))$
- $f \overline{\oplus} G = (wz') \overline{\oplus} ((c \oplus (za + wa'))'(d \oplus (zb + wb')))$
- $M(a,b,c,d) = ab'c'd' + a'bcd$

Extended matching

- Captures implicitly all possible matches.
- No extra burden when exploiting *don't care* sets.
  - $M(\mathbf{c}) = \forall_{\mathbf{x}} [G(\mathbf{x}, \mathbf{c}) \; \overline{\oplus} \; f(\mathbf{x}) + f_{DC}(\mathbf{x})]$
- Efficient BDD-based representation.
- Extensions to support multiple-output matching

Summary

- Library binding is very important.
- Rule-based approach:
  - General, sometimes inefficient.
- Algorithmic approach:
  - Pattern-based: fast, but limited.
  - Boolean: more general and efficient.